

# On Adaptive HTTP Streaming to Mobile Devices

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# Outline

- ❑ Motivation & Background
- ❑ System Design
- ❑ Adaptation and Battery Consideration
- ❑ Performance Evaluation
- ❑ Conclusion

# DASH over Mobile Devices

## ❑ DASH: Dynamic Adaptive Streaming over HTTP

- single video streamed at multiple versions
  - address user heterogeneity: access/device
- existing infrastructure optimized for HTTP
  - CDN, Caching, Firewall, etc

## ❑ DASH over Mobile Devices

- wide adoption of mobile devices
  - smart phones, tablets
- more limitations on mobile devices
  - wireless connection
    - Wifi, cellular
  - battery limitation

What?  
Battery  
Out!!!!



# Video Adaptation

## ❑ Which Version to Download?

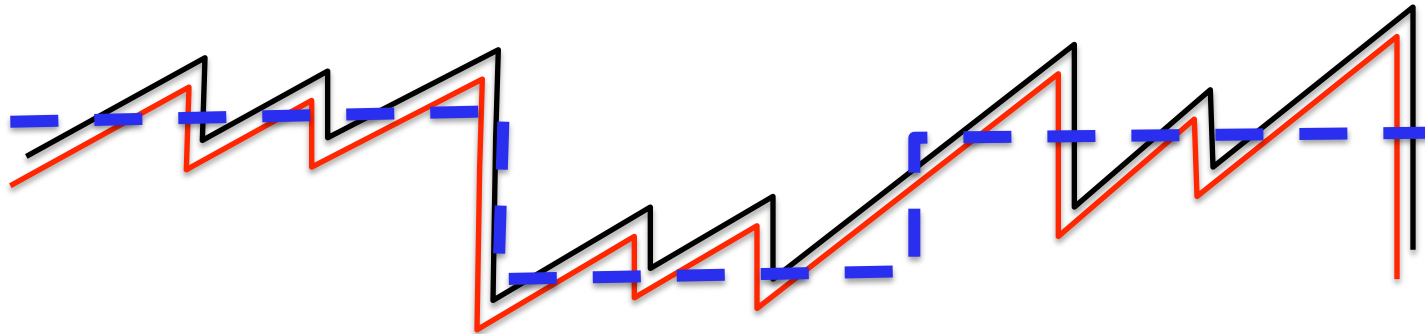
- download rate controlled by TCP
- version selection: application layer control
- video rate should match TCP throughput
  - too aggressive: video freezes,
  - too conservative: low video quality

## ❑ TCP throughput prediction

- TCP throughput has variations at different time scales

## ❑ Video rate smoothing

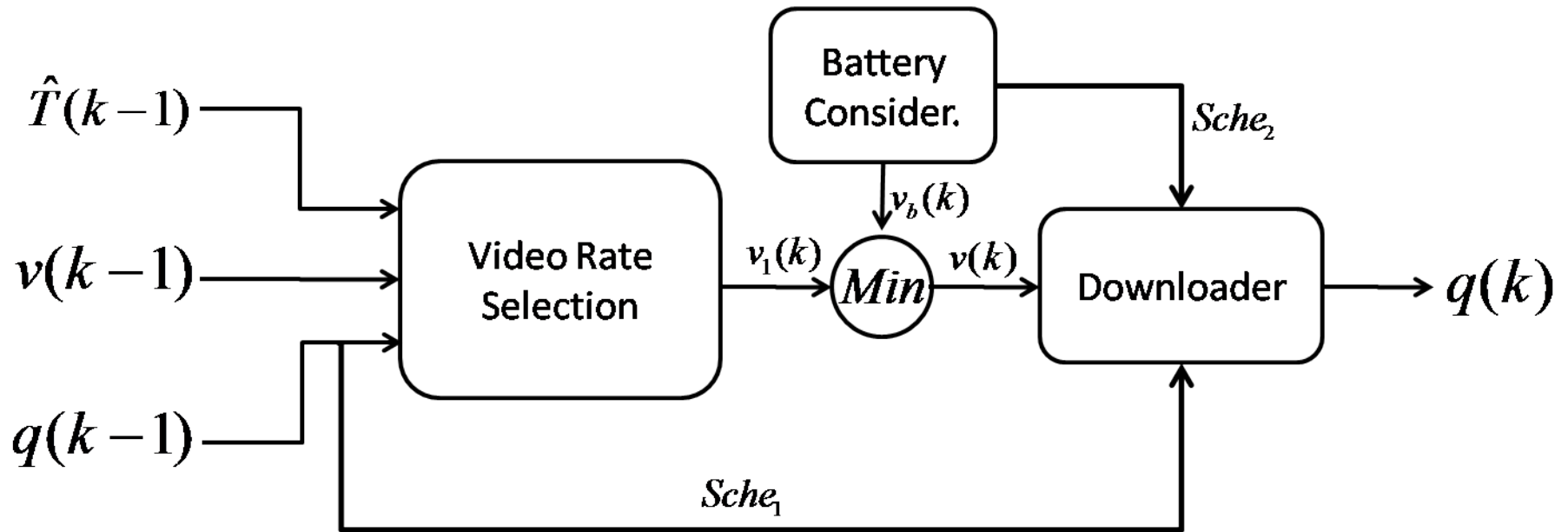
- users sensitive to frequent quality changes



# Challenges

- ❑ Variability of TCP throughput
  - variability: built-in character of TCP
  - network congestion and server overload
  - wireless connection only on mobile devices
- ❑ Management of video playback buffer
  - small buffer: Low delay, high risk of freezing
  - large buffer: Safe but waste of resources (especially for cellular ), impossible for live streaming
- ❑ Battery limitation on mobile devices
  - different energy consumption patterns on WiFi and cellular
  - what video rate should be chosen when battery is depleting?

# Design Framework



$v_1(k)$ : target rate based on buffer size, TCP throughput, previous rate

$v_b(k)$ : target rate based on battery consideration

$v(k)$ : requested rate  $\text{Min}(v_1(k), v_b(k))$

# TCP throughput Based Adaptation

- ❑ Simple History-Based prediction( $\hat{T}(k)$ )
  - Proved to generate highly accurate prediction given real time throughput data
  - $N = 10$  (Number of historical data)
- ❑ Dynamic Margin ( $M$ )
  - target video rate lower than throughput estimate
  - rate margin depends on TCP throughput variability
    - higher variability  $\rightarrow$  larger margin

$$SI(k) = \frac{1}{N-1} \sum_{i=k-N+1}^{k-1} \frac{|T(i) - T(i-1)|}{T(i)} \quad M(k) = 1 - \frac{25 + 70e^{SI(k)}}{100e^{SI(k)}}$$

- ❑  $\hat{v}_1(k) = Q(\hat{T}(k)(1 - M(k)))$  for smoothing

# Smooth Video Adaptation Algorithm

## □ Prompt Rate Decrease

- if buffer becomes lower than a threshold
- avoid buffer underflow (freeze)

## □ Conservative Rate Increase

- only if target rate higher than the current rate for  $m$  consecutive steps
- $m$  calculated dynamically
- avoid oscillations triggered by transient TCP rate increase

## □ Control Buffer Overflow

- introduce idle time between chunk requests

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### Algorithm 1 Smooth Video Adaptation Algorithm.

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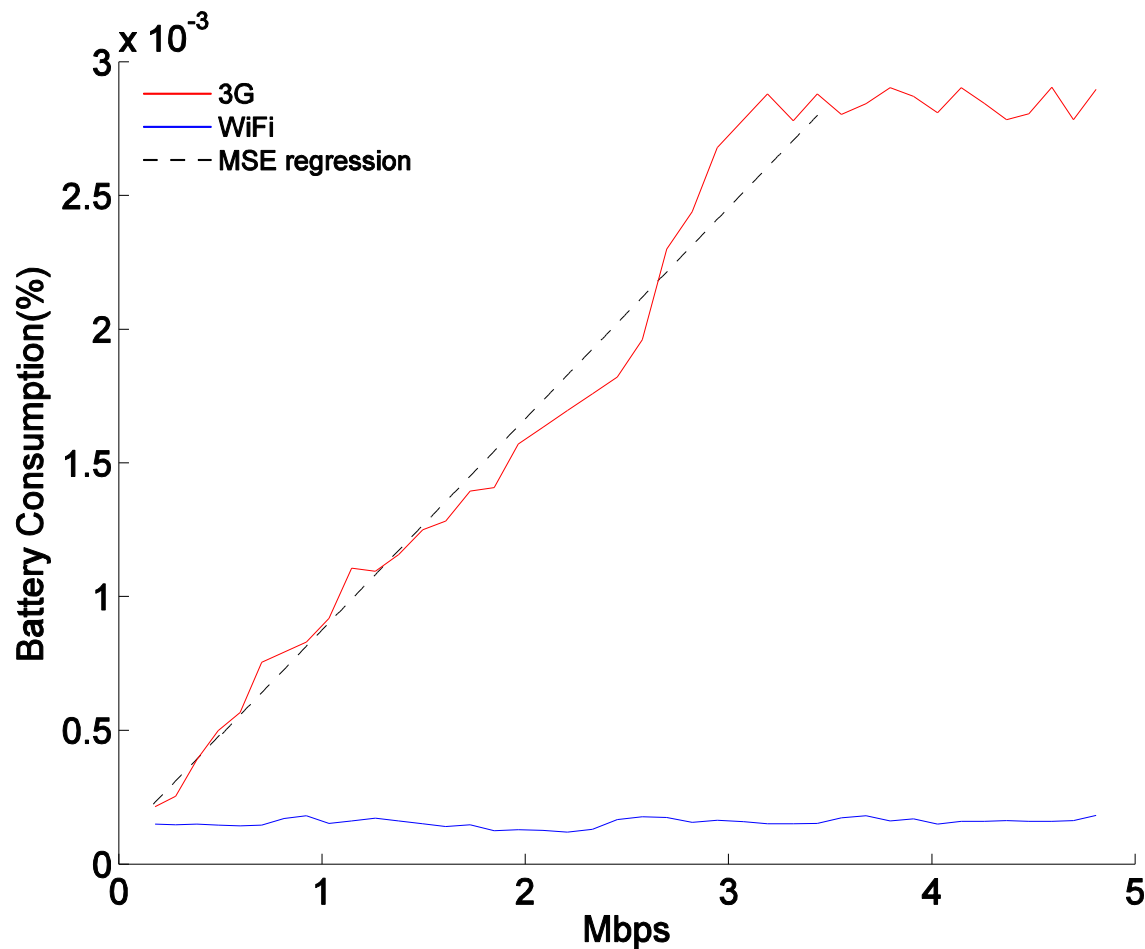
```
1: if  $q(k) < \frac{q_{th}}{2}$  then
2:    $v(k) = Q(\bar{T}(k-1)(1 - M(k)))$ ;
3:   return;
4: else
5:    $\hat{v}(k) = Q(\hat{T}(k)(1 - M(k)))$ ;
6:   if  $\hat{v}(k) \geq v(k-1)$  then
7:     Counter ++
8:     if Counter >  $m$  then
9:        $v(k) = \hat{v}(k)$ ;
10:      Counter = 0;
11:      return;
12:   end if
13: else
14:   Counter = 0
15: end if
16: end if
17:  $v(k) = v(k-1)$ ;
18: if  $q(k) - q_{cap} > 0$  then
19:    $I = q(k) - q_{cap}$ ;
20: else
21:    $I = 0$ 
22: end if
23: Idle( $I$ );
24: return;
```

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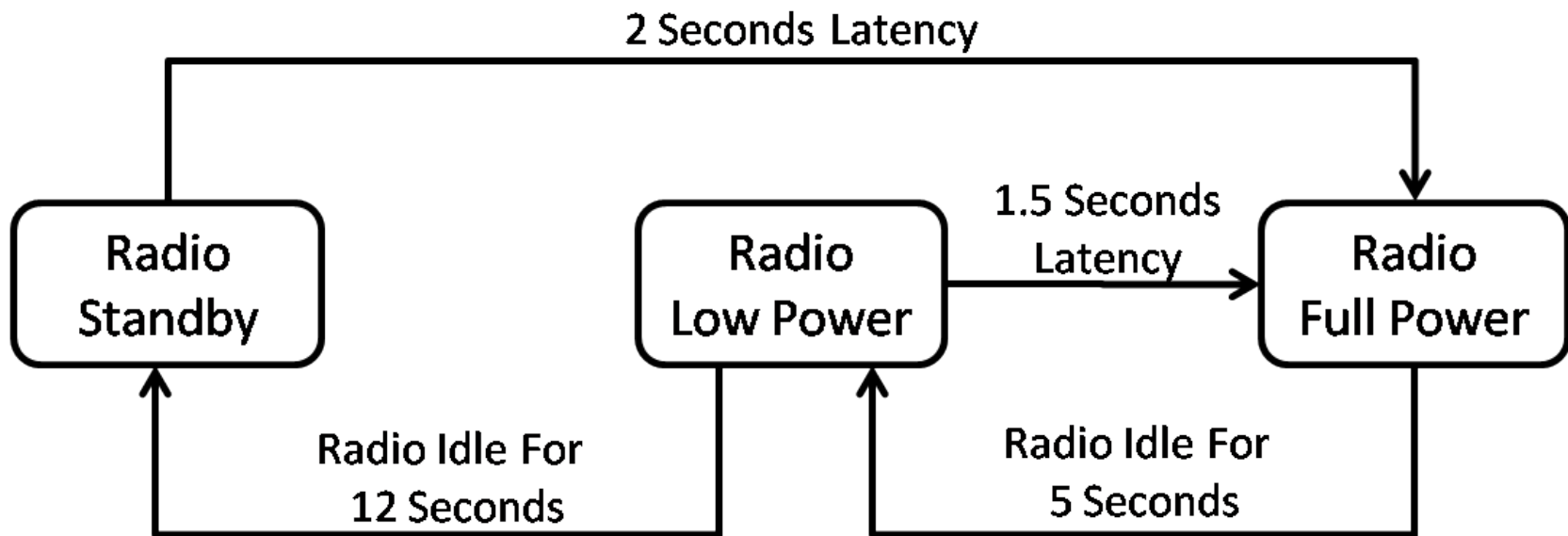
# Battery Consideration

- ❑ 3G radio consumes more battery than WIFI



# Battery Consideration

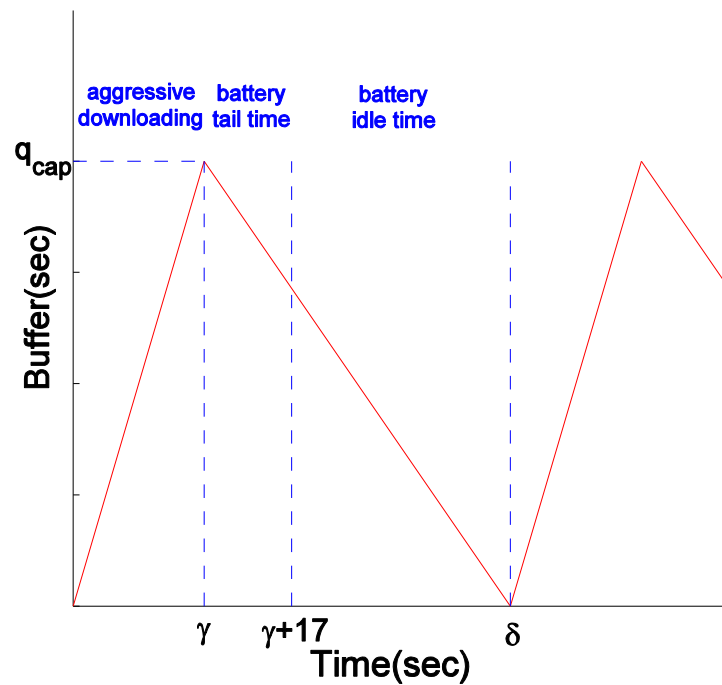
## ❑ 3G Battery State Machine (AT&T)



Different Scheduling can make huge difference

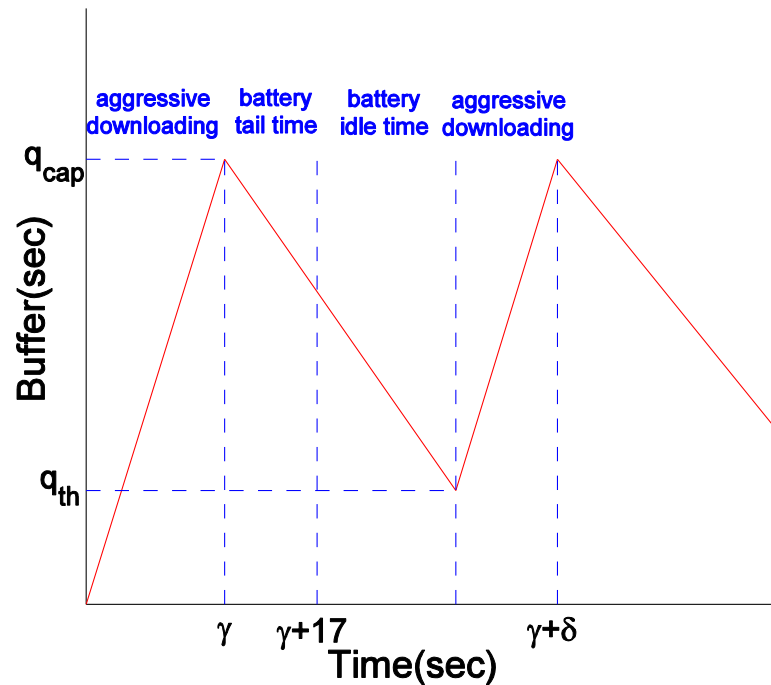
# Scheduling Considering 3G Battery mode

- Download more chunks in one transaction given buffer size bounds



Ideal case

$$\gamma = q_{cap} * \frac{v}{T - v} \quad \text{: Accumulating from 0 to cap}$$



Practical case

$$\sigma = q_{cap} * \frac{T}{T - v} \quad \text{: Loop Period}$$

## Battery Life Constraint

- Get the highest rate at which the batter can last for the whole movie

$$v_b(r, p) = \max(v_i, v_i : \frac{p}{E_{v_i}} > r)$$

$E_{v_i}$  : Unit Time Energy Consumption for rate  $v_i$

$r$  : Remaining video time

$p$  : Remaining battery percentage

- $E_{v_i}$  calculated by MSE regression of 3G battery consumption data

- $v(k) = \min(v_1(k), v_b(k))$

# Evaluation: test platforms

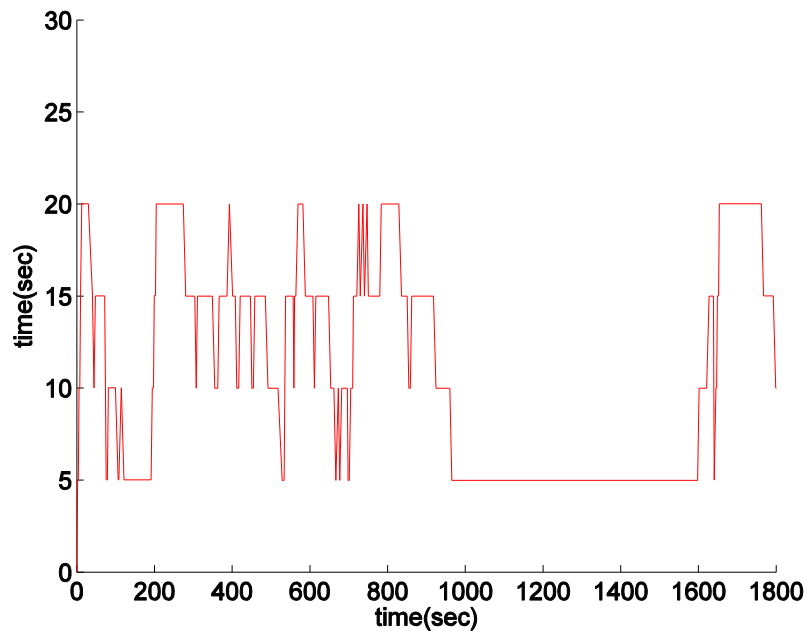


- ❑ Vanilla apache HTTP server, Google Nexus 4 with Andriod 4.2 as client
- ❑ Video Rates
  - 100Kbps~4.1Mbps with even rate gap of 100Kbps
  - five-second video chunks
- ❑ Internet Experiments
  - Wifi connection: change the distance between laptop and router
  - 3G connection: in motion
  - Hand over between 3G and Wifi

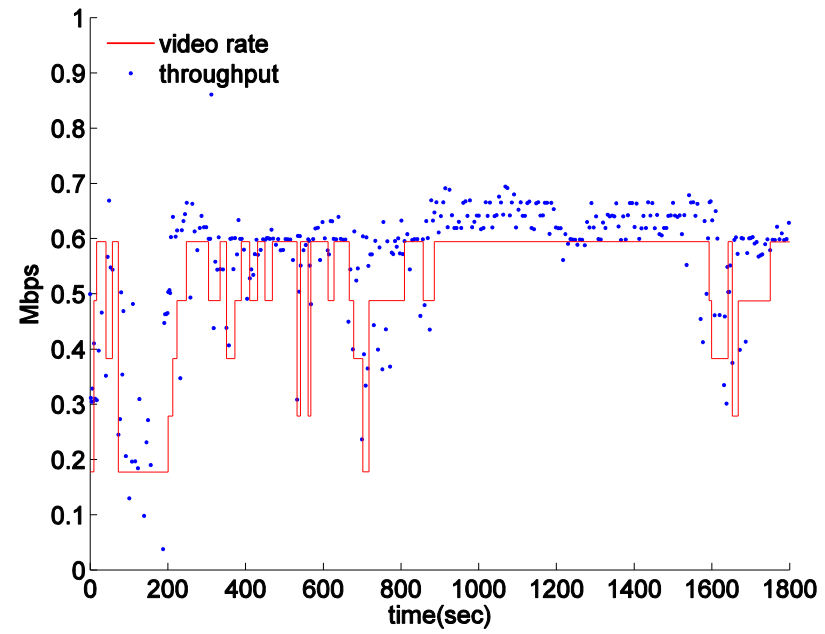
# TCP Throughput Variability

(Mbps)	Wired	Strong WiFi	Weak WiFi	Static 3G	3G In motion
SI	0.004	0.074	0.336	0.151	0.214
Average Throughput	0.720	0.699	0.616	2.597	2.488
Variation	0.001	0.002	0.129	0.164	0.710

## Strong WiFi without margin, $q_{cap}=20$



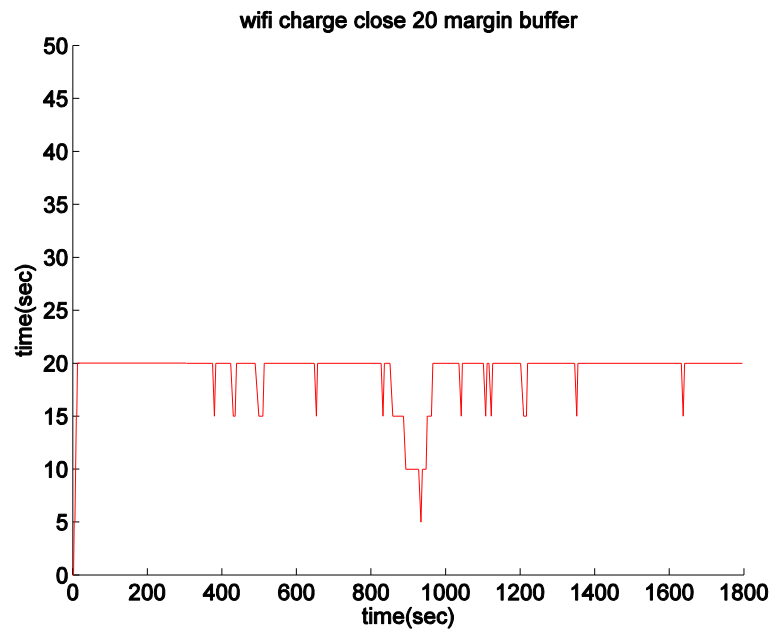
video buffer length



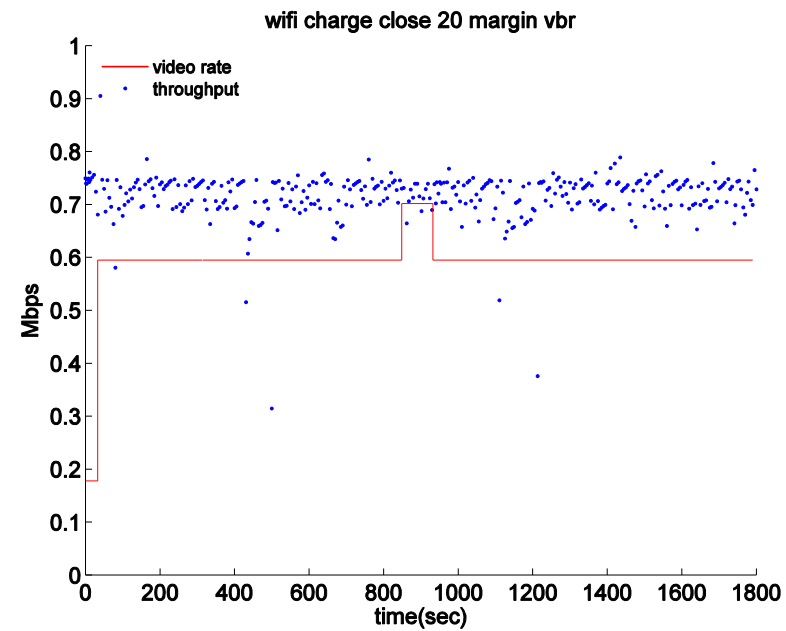
video rate vs. TCP throughput

- ❑ Video rate follows TCP throughput tightly, but too much oscillations for users.
- ❑ Buffer fluctuates a lot
- ❑ Bad User QoE

# Strong WiFi with margin, $q_{cap}=20$



video buffer length

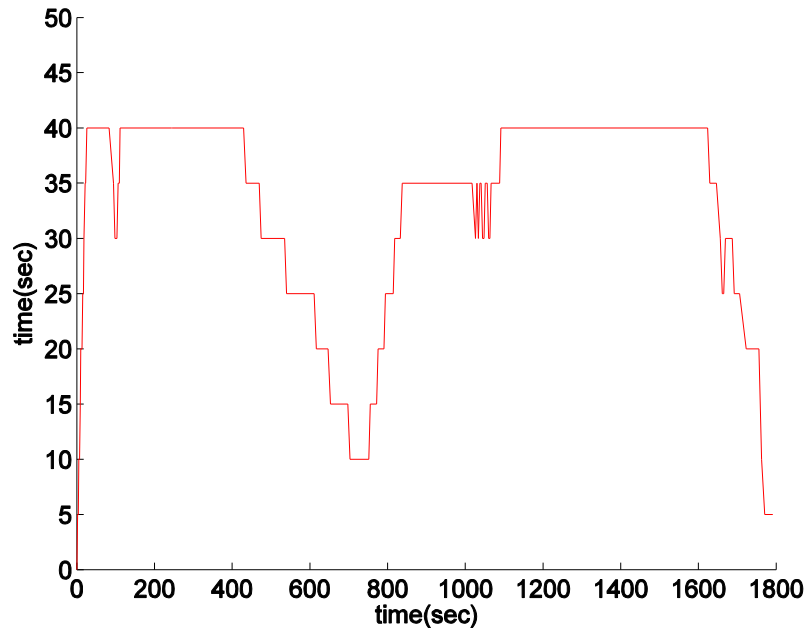


video rate vs. TCP throughput

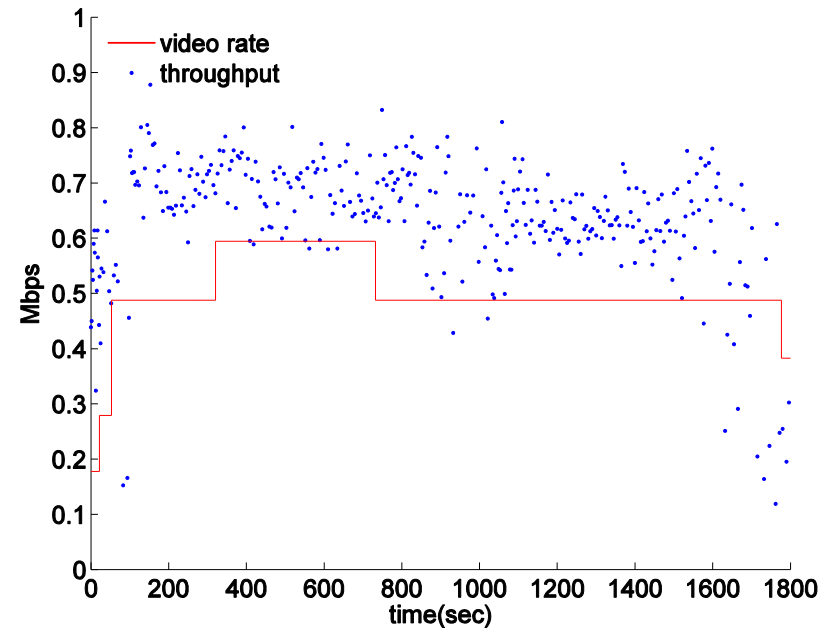
- ❑ Video rate slightly lower than throughput
- ❑ Stable buffer and video rate
- ❑ Good user QoE



## Weak WiFi with margin, $q_{\text{cap}} = 40$



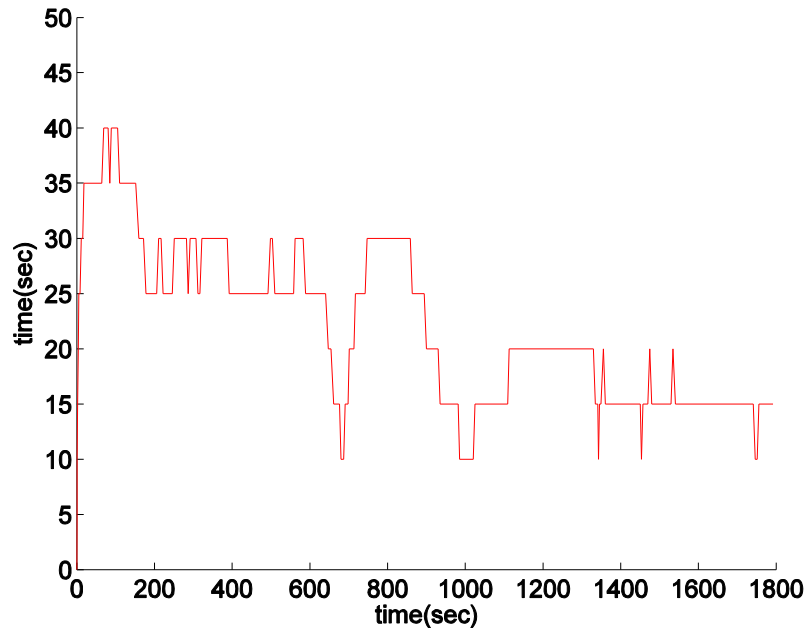
video buffer length



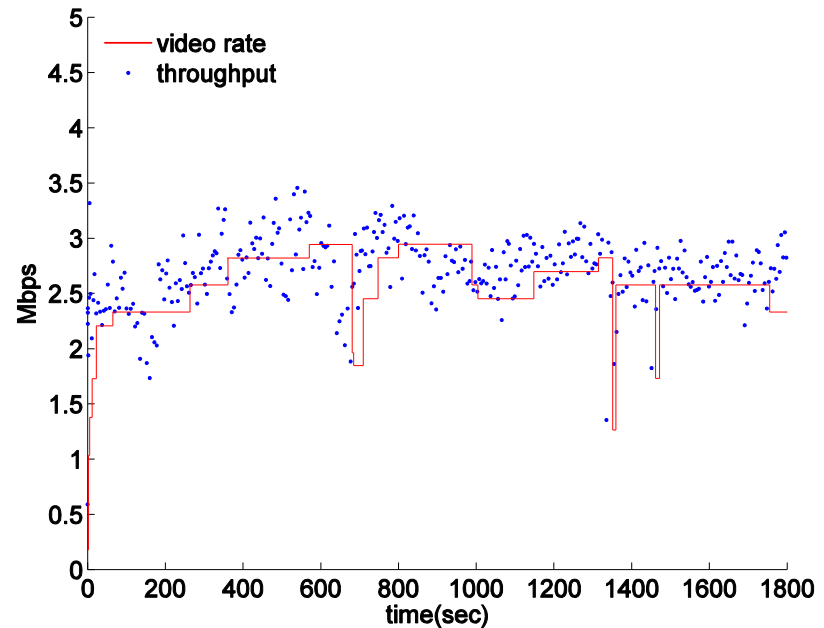
video rate vs. TCP throughput

- ❑ Buffer fluctuates but never hit zero (transparent to users)
- ❑ Stable video rate, good user QoE
- ❑ Lower buffer cap value won't work

## 3G without margin, $q_{cap}=40$ , $q_{th}=20$



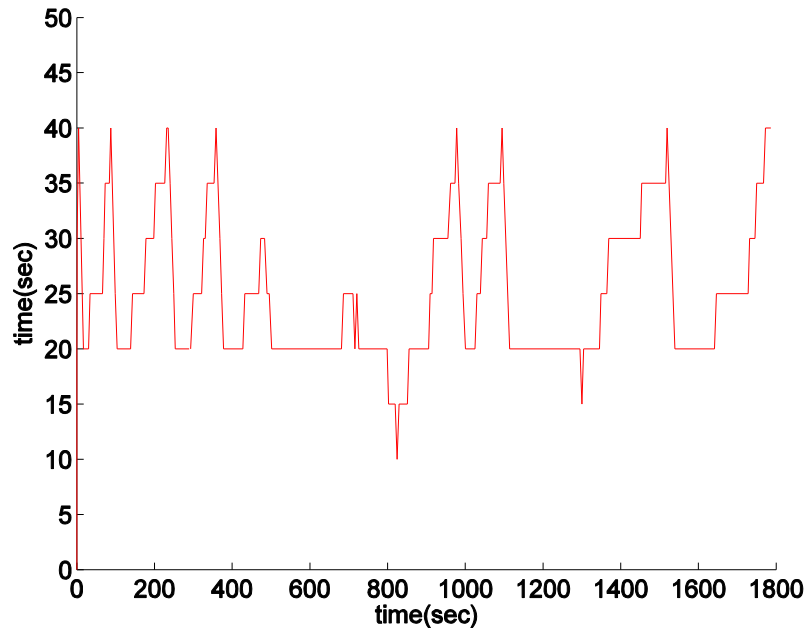
video buffer length



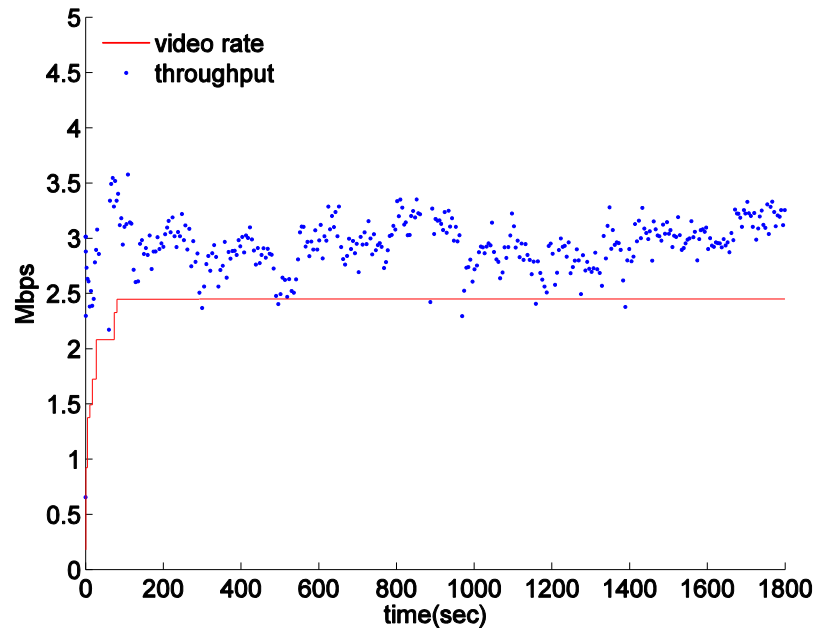
video rate vs. TCP throughput

- ☐ Buffer and video rate fluctuates
- ☐ Bad user QoE
- ☐ Pretty high throughput

## 3G with margin, $q_{\text{cap}} = 40$ , $q_{\text{th}} = 20$



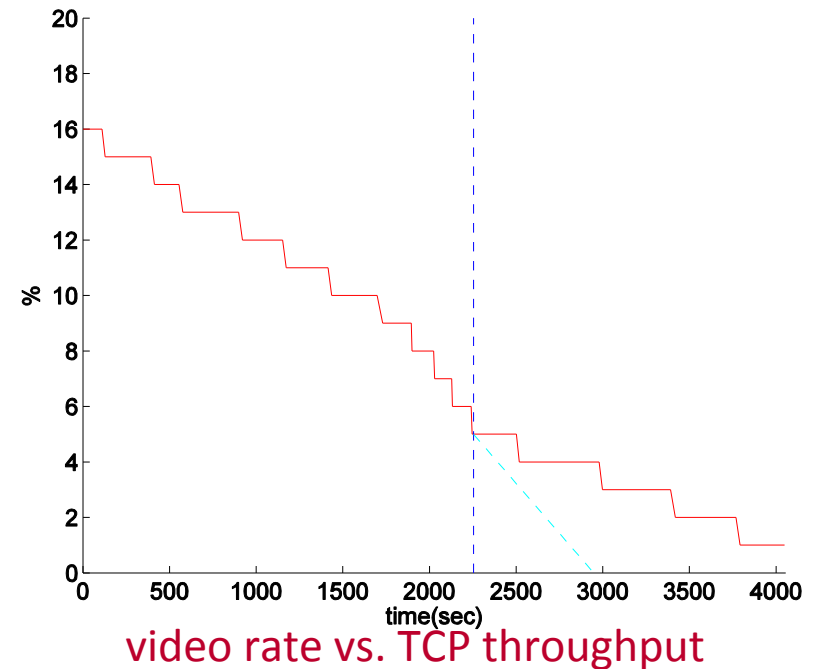
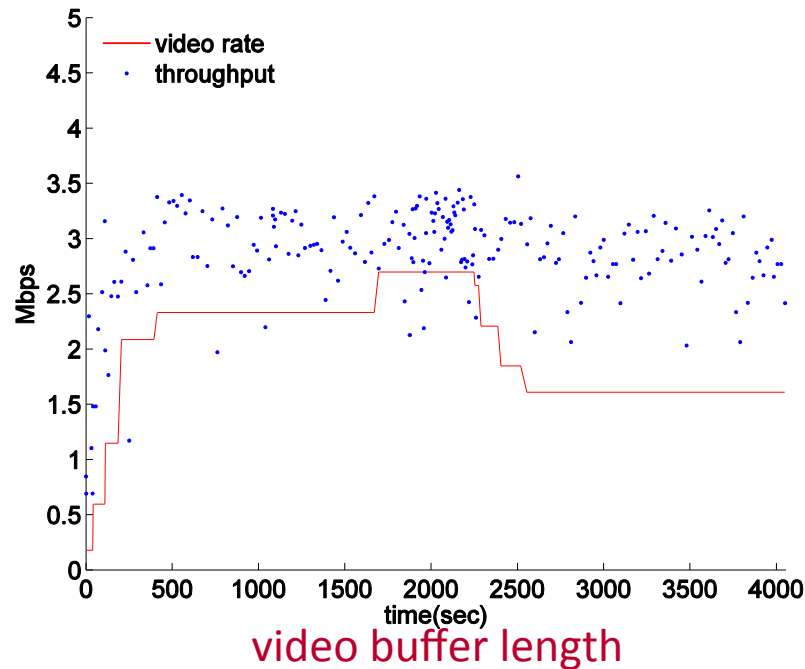
video buffer length



video rate vs. TCP throughput

- ☐ Buffer fluctuates as scheduled through bundled chunk downloading mechanism
- ☐ Stable Video rate
- ☐ Good user QoE

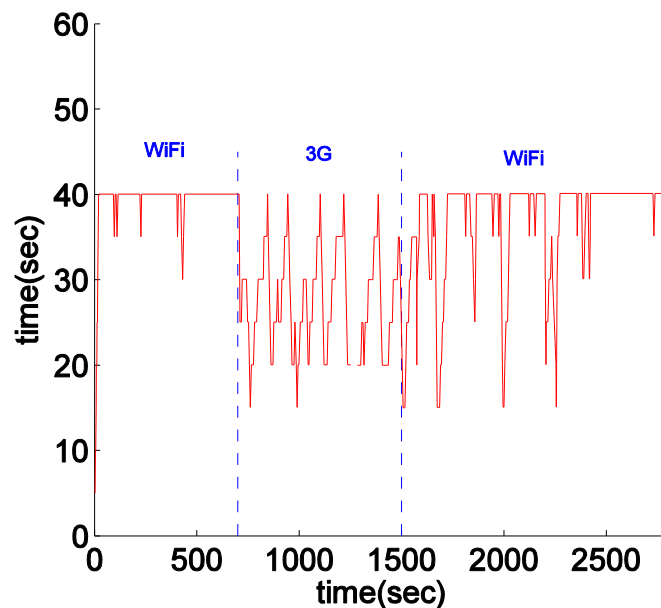
# Battery Consumption and Video Rate Adaptation over 3G Connection



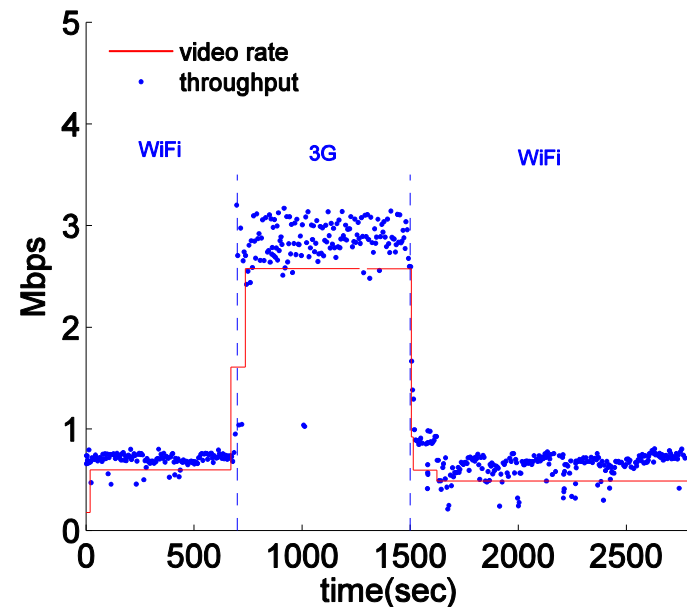
- ❑ Video rate drops because of battery limitation while throughput still high
- ❑ Battery consumption rate decreased to finish the whole video

# Handover between WiFi and 3G connection

$$q_{\text{cap}}=40, q_{\text{th}}=20$$



video buffer length

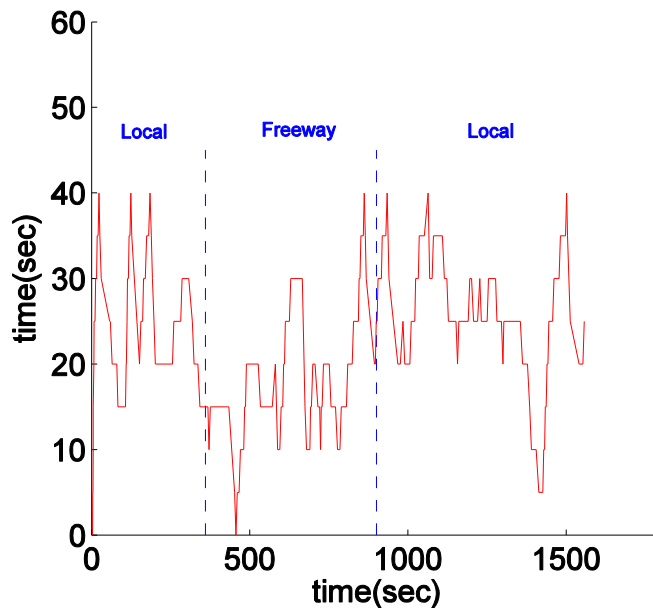


video rate vs. TCP throughput

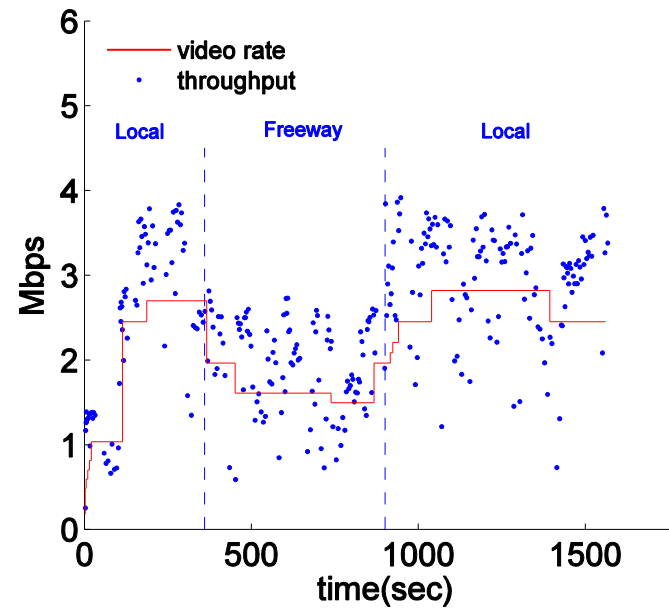
- ☐ Buffer fluctuates as scheduled in 3G and stable in WiFi
- ☐ Stable Video rate with responsive handover
- ☐ Good user QoE

# 3G Mobile Experiments in a Moving Car

$$q_{\text{cap}} = 40, q_{\text{th}} = 20$$



video buffer length



video rate vs. TCP throughput

- ❑ Buffer fluctuates a lot (Hits zero once!)
- ❑ Highly fluctuating throughput at high speed (30~70mph)
- ❑ But still acceptable QoE

# Conclusion

- ❑ Client-side video adaptation algorithm
  - TCP throughput highly dynamic
  - Dynamic rate margin leads to smooth video rate
  - Video scheduling with battery considerations (3G on-off pattern and available battery level)
  - Extensive evaluation on real systems
- ❑ Future Work
  - Considering cooperation between mobile devices

**Thanks!**

**Q&A**