

Packet Video Workshop

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# OPTIMAL ADAPTATION TRAJECTORIES FOR BLOCK-REQUEST ADAPTIVE VIDEO STREAMING

# Outline

- ① Motivation
- ② Problem formulation
- ③ Approach
- ④ Evaluation

# Adaptive Video Streaming

- ◎ Cisco forecast: video will be 80 % - 90 % of total consumer traffic in 2017
- ◎ Huge heterogeneity of terminal devices
  - Screen sizes, screen resolutions, CPU power, battery capacity
- ◎ Huge heterogeneity of network conditions
  - Throughput from 10 kbps to 10 Mbps and more
  - Packet loss rate from 0 to 10 % and more
  - Latency from 1 ms to 1 s and more

Adaptation necessary

# Adaptive Video Streaming

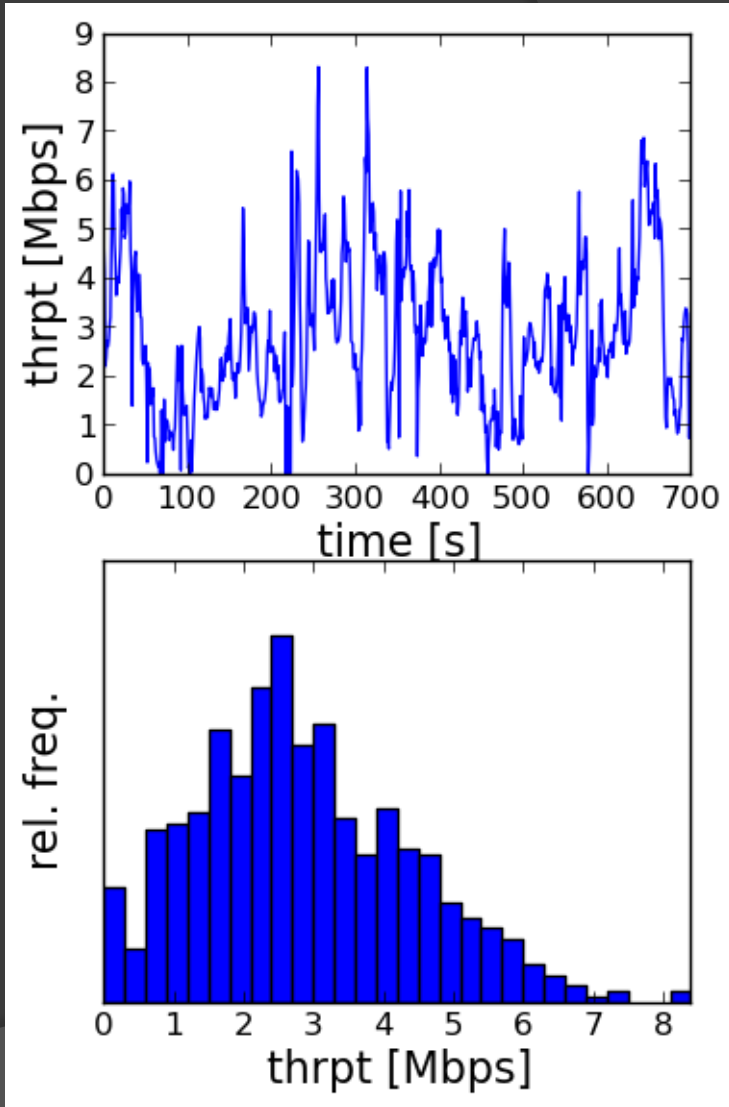
- ⦿ Network conditions change dynamically
  - Cross-traffic
  - Mobility
    - Outdoor/indoor, overground/underground
    - Changing distance to base station
    - Changing environment geometry
- ⦿ Continuous dynamic adaptation required, e.g.,
  - Adaptation of video bit-rate
  - Adaptation of other encoding parameters
    - Example: GOP size to loss rate
    - ...

# State-of-the-Art

- ◎ Block-request adaptive streaming
  - Stateless server, e.g., HTTP
  - Client requests chunks of video data
  - Chunk representation selected dynamically
    - Based on network conditions, etc.
- ◎ Standards: MPEG-DASH, HLS (draft), etc.
- ◎ Popular commercial implementations
  - Microsoft SmoothStreaming (proprietary)
  - Apple QuickTime, iOS (HLS)
  - Adobe HTTP Dynamic Streaming (proprietary)

# Adaptation strategies

- Adaptation strategy is essential for QoE
- Challenge: random throughput, high variance (esp. wireless networks)
- Conflicting objectives
  - Avoid underruns
  - Max. average media bit-rate
  - Min. quality jumps
  - Min. start-up delay



# Problem: How to evaluate adaptation strategies?

- ⦿ Compare to predefined requirements
  - No predefined requirements: best-effort
  - User's expectations depend on many factors
    - Age, affinity to technology, viewing context, etc.
- ⦿ Compare to state-of-the-art solution
  - No widely accepted state-of-the-art clients/benchmarks
- ⦿ Compare to optimum
  - Calculate optimal adaptation trajectories for given network conditions

# Approach



# Approach and application

- ⊙ Calculate optimal adaptation trajectories, given
  - Throughput over time
  - Segment size and representation information
- ⊙ Application scenario I
  - Record clients throughput over time
  - Calculate optimum and compare
  - Difficulty: client might introduce delays between requests  
→ Potential loss of optimality
- ⊙ Application scenario II
  - Rerun continuous TCP flow under same conditions
  - Calculate optimum and compare
  - Multiple runs to account for randomness

# Further applications

- ⦿ Evaluate influence of various factors on achievable performance
- ⦿ Influence of video parameters
  - Number and bit-rates of representations
  - Segment duration
  - ...
- ⦿ Influence of network parameters
  - QoS: Throughput, packet loss, latency
  - MAC strategy
  - TCP flavor
  - ...

# Optimality metric

- ⦿ Ultimate goal: optimize QoE
  - ⦿ Factors that influence QoE
    - Re-buffering duration and distribution over time
    - Average quality over chunks, minimum quality
    - Number of quality switches and distrib. over time
    - Start-up delay
  - ⦿ No unifying QoE metric exists so far
  - ⦿ Our optimization objectives and constraints
    - No re-buffering
    - Maximum average video bit-rate
    - Minimum number of switches
    - Start-up delay is configuration parameter
- } two step approach

# Step 1: Average bit-rate maximization

- Given
  - Throughput over time  $V(t)$
  - Desired start-up delay
- Objective: maximize average video bit-rate
- Additional constraint: no buffer underruns

$$\begin{aligned} \text{(OP1)} \quad & \max \quad \sum_{i=1}^n \sum_{j=1}^m S_{ij} x_{ij} \\ & \text{s. t.} \quad \sum_{j=1}^m x_{ij} \geq 1 \quad \text{for all } i = 1, \dots, n \\ & \quad \quad \sum_{i=1}^k \sum_{j=1}^m S_{ij} x_{ij} \leq V(D_k) \quad \text{for all } k = 1, \dots, n \end{aligned}$$

$i = 1, \dots, n$  - segments  
 $j = 1, \dots, m$  - representations  
 $S_{ij}$  - segment size  
 $D_i$  - playback deadline  
 $x_{ij}$  - download  $i$  from  $j$

- MCNKP
- NP-hard
- Pseudo-polyn.
- Solution: seconds (Gurobi)

# Step 2: Quality switches minimization

- Given: as in step 1 plus optimal average quality  $V^*$
- Objective: minimize number of switches
- Constraint: same as step 1 plus average quality equal to optimum

$i = 1, \dots, n$  - segments  
 $j = 1, \dots, m$  - representations  
 $S_{ij}$  - segment size  
 $D_i$  - playback deadline  
 $x_{ij}$  - download  $i$  from  $j$   
 $V^*$  - optimum val. of (OP1)

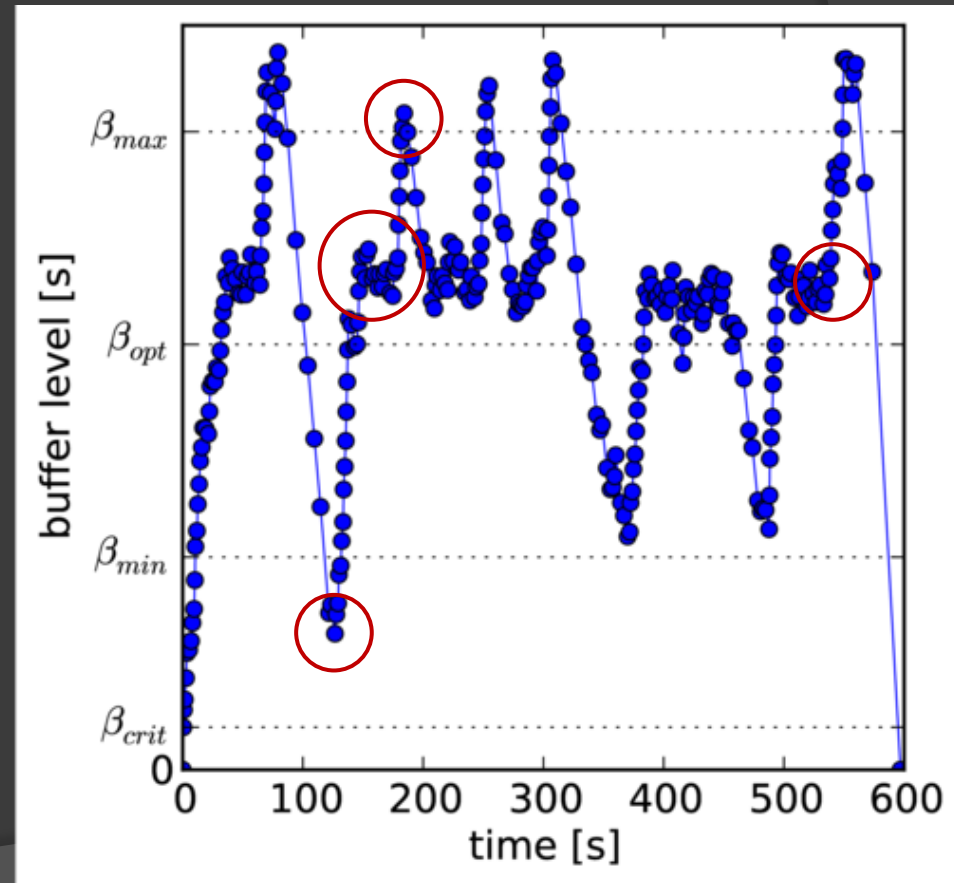
$$\begin{aligned}
 \text{(OP2)} \quad & \min \quad \frac{1}{2} \sum_{i=1}^{n-1} \sum_{j=1}^m (x_{ij} - x_{i+1,j})^2 \\
 \text{s. t.} \quad & \sum_{j=1}^m x_{ij} \geq 1 \quad \text{for all } i = 1, \dots, n \\
 & \sum_{i=1}^k \sum_{j=1}^m S_{ij} x_{ij} \leq V(D_k) \quad \text{for all } k = 1, \dots, n \\
 & \sum_{i=1}^n \sum_{j=1}^m S_{ij} x_{ij} \geq V^*
 \end{aligned}$$

- Quadratic MCNKP
- Solution: minutes (Gurobi)

# Evaluation

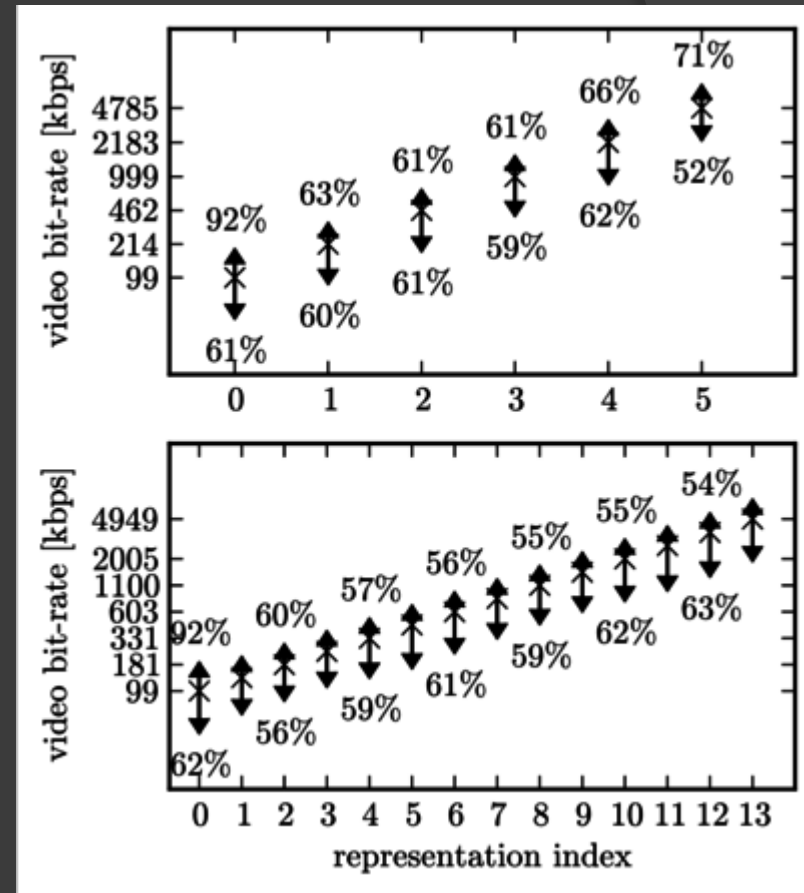
# Previous work: adaptation strategy

- ⦿ DASH does not specify adaptation strategy
- ⦿ Developed own algorithm (impl. as plugin for VLC)
- ⦿  $\beta < \beta_{min} \wedge \beta' < 0 \Rightarrow \searrow$
- ⦿  $\beta > \beta_{max} \wedge \beta' > 0$ 
  - $r^\uparrow < \alpha\rho \Rightarrow \nearrow$
  - $r^\uparrow \geq \alpha\rho \Rightarrow \rightsquigarrow$
- ⦿  $\beta > \beta_{opt} \wedge r^\uparrow \geq \alpha\rho \Rightarrow \rightsquigarrow$
- ⦿  $\beta < \beta_{crit} \Rightarrow \Downarrow$
- ⦿ + aggressive at start-up
- ⦿ + some additional tweaks
- ⦿ In total: 10 parameters



# Video used for evaluation

- Big Buck Bunny (animated)
  - 598 seconds
- Encoded in 6 and 14 representations
  - 299 segments, 2 sec. each
  - Bit-rates logarithmically from 100 kbps to 5 Mbps
  - Kept bit-rates fluctuations low
  - 2 manifests and container formats: DASH and MSS
- Low fluctuation amplitude is important if segment size not known in advance

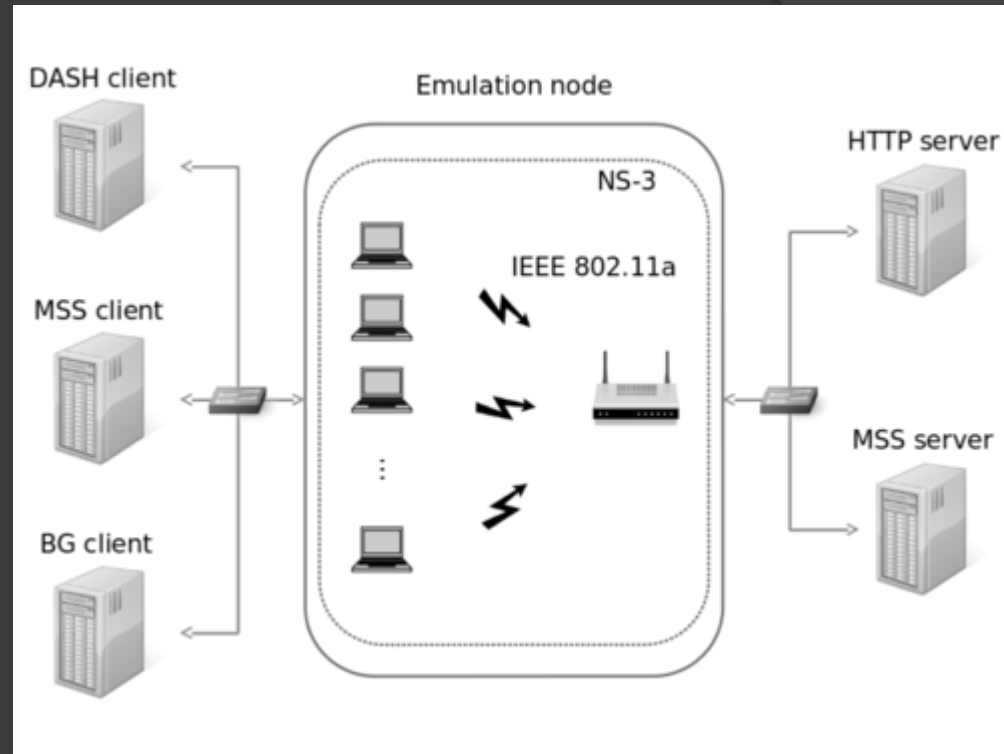


(Note the log y-axis)



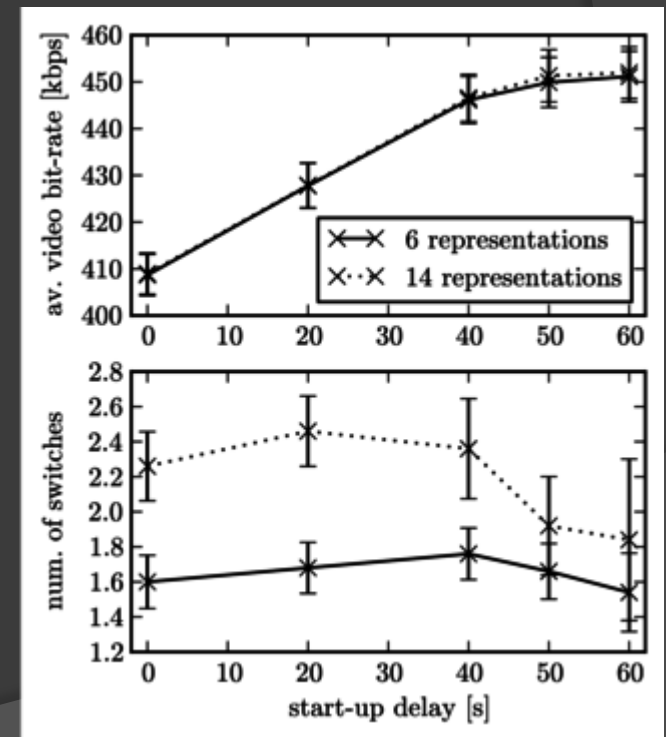
# Evaluation setting

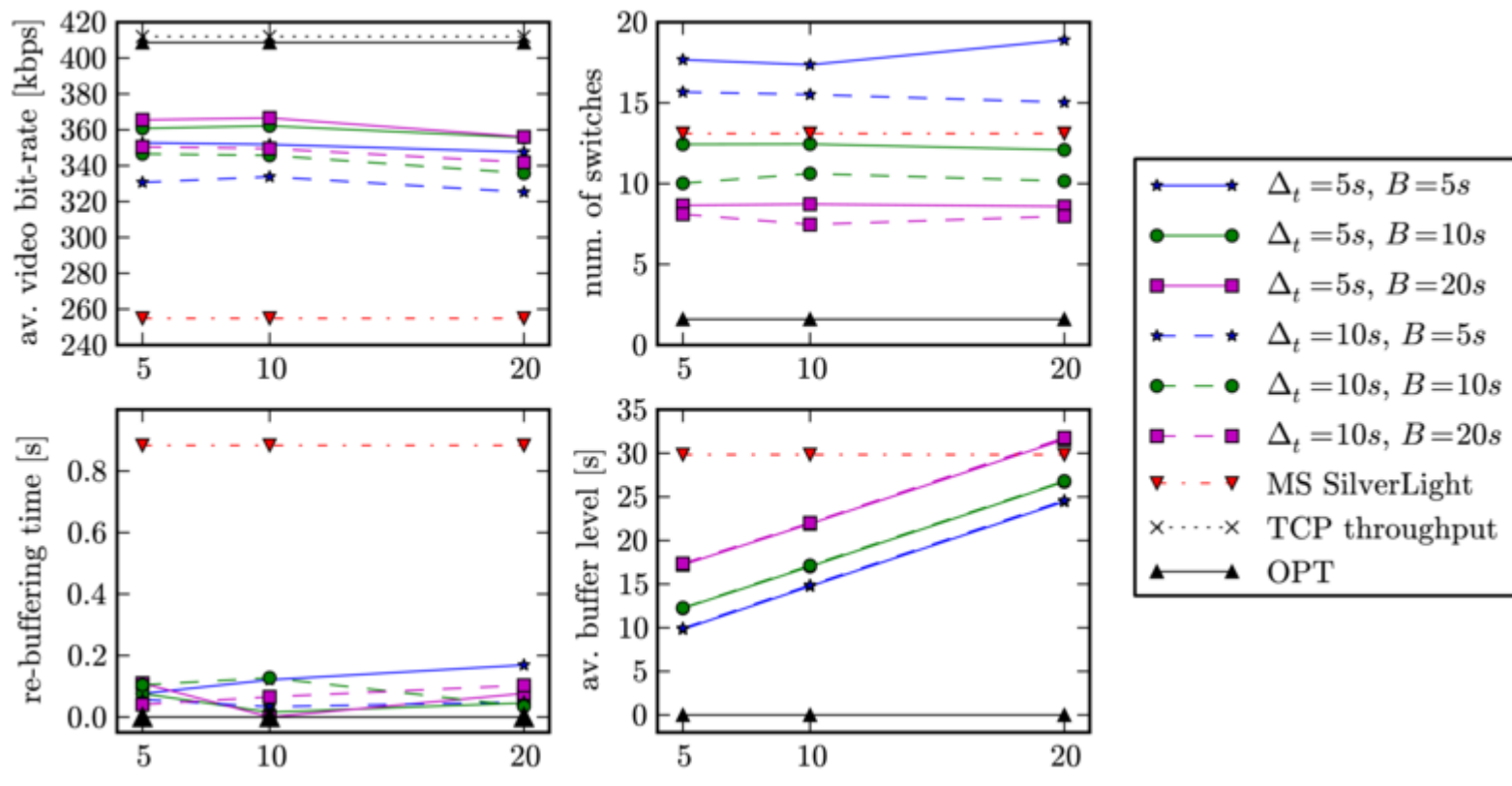
- 802.11a model based on BOWL indoor testbed (7 stations)
- Max. TCP throughput: 1.4, 1.7, 19, 19, 21, 21, 21 Mbps
- Second slowest selected for video traffic
- 14 synthetic HTTP clients as cross-traffic (Pries et al.)
  - 2 on each wireless station
  - Detailed model: experimentally fitted distributions for
    - User activity, main object sizes, secondary object sizes, inter-object intervals, etc.



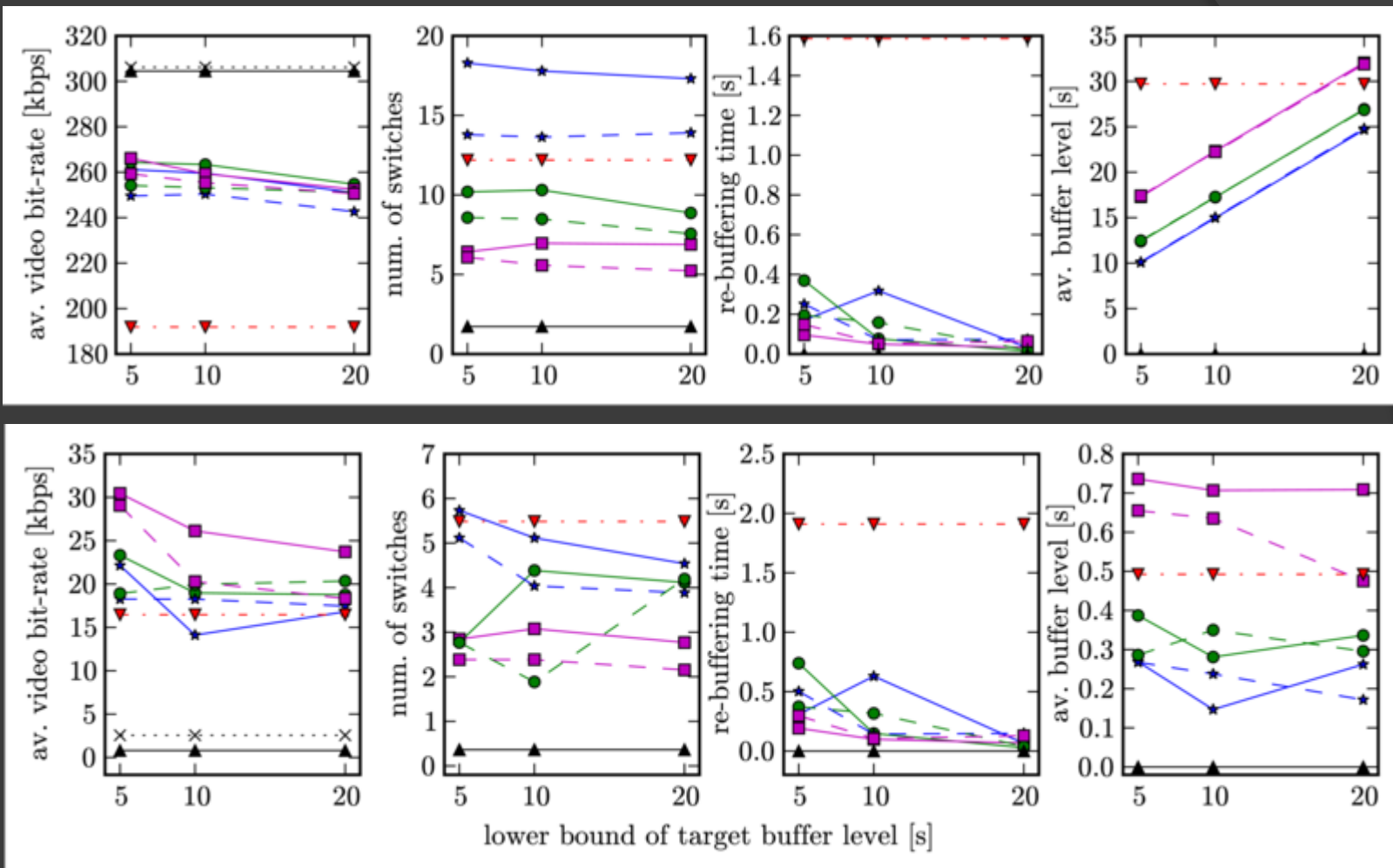
# Results: influence of video parameters

- Influence of number of video representations
- Influence of start-up delay
- 6 representations sufficient
- Num. switches surprisingly low (here: upper bounds)
- Start-up delay has little influence
  - Approx. 12 % for 60 s





- OPT utilize almost 100% of TCP's fair share
- OPT has 0 re-buffering, very little switches, almost 0 buffer level
- DASH av. video bit-rate: 78% to 90% of OPT, MSS: 62%
- Num. of switches, re-buffering: DASH is better or comparable
- DASH has lower avg. buffer level → better for live content
- Good DASH configuration:  $\beta_{min} = 10s, B = 20s, \Delta_t = 5s$



- Two clients on same wireless stations (backgr. as before)
- Differences between two clients, averaged over runs
- Good fairness w.r.t. avg. bit-rate, re-buffering, buffer level
- Medium fairness w.r.t. number of switches

# Conclusion

- ⊙ Optimal adaptation trajectories allow to
  - Benchmark adaptation strategies
  - Study influence of network and video parameters
- ⊙ Potential extensions
  - Optimize w.r.t. QoE metric, once available
- ⊙ Evaluation
  - DASH comparable or better in studied setting
  - DASH achieves 78% to 90% of optimum

THE END

# References

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