## Are Information-Centric Networks Video-Ready?

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#### Presentation

- Discussion
  - No specific solution/system presented
  - Highlight good and not so good features of ICN w.r.t.
    video transport
  - Point issues that need attention
- Why is this discussion important?
  - Video applications attracted ICN researchers
  - Prototype implementations focus on message passing using ICN primitives
  - Critical aspects w.r.t. performance and scalability left for future work

### Our experience with ICN

• Participated in



2008 - 2010



2011 - 2013

- Publish-Subscribe Internetworking (PSI)
- Implemented video applications in prototypes [1]
  - Appealing demos
  - Promising application
- Message passing but not deep study of application behavior
  - Many core pieces of the network architecture still missing
- [1] Parisis et al., "Demonstrating Usage Diversity Over an Information-Centric Network," demo in IEEE INFOCOM 2013.

#### Can we finalize some aspects on ICN and move on?

- Many ICN proposals
  - Content-Centric Networking
  - NetInf
  - Publish-Subscribe Internetworking
  - ...
- With similarities
  - Goal: Primary focus to content distribution
  - Self-identified information items
  - Universal caching, anycast, multicast
- And differences
  - Diverse approaches in core functions
    - Item lookup, routing, forwarding
  - CCN: pull-based, distributed control plane, hop by hop routing/forwarding
  - PSI: push-based, centralized control plane, explicit-routing

### Internet Video Transfer

- Internet video applications operate on top of <u>well defined</u> <u>architecture</u>
  - End-to-end system design
  - Network layer: best effort, IP host addresses
  - Transport layer: UDP, TCP
  - Application layer: RTP, HTTP
- Applications choose protocols based on application context and protocol behavior
  - Video on Demand vs Live Streaming
  - Stream adaptation
- Can we port existing video applications to ICN as is?
  - ICN API looks similar to application layer protocols...
- Not that simple

### Rest of presentation

- Two diverse ICN architectures
  - Content-Centric Networking
  - Publish-Subscibe Internet
- Two kinds of video applications with different transport requirements
  - Video on Demand: reliable transfer
  - Live Streaming: real-time delivery
- Which features of ICN facilitate video transfer
- What seems problematic

### Content-Centric Networking (CCN)

- Named content packets
  - Hierarchical names
  - Interest Data packets
  - No host addresses
- Pull-based operation
  - One Interest per Data
- Packet caches in routers
- Native multicast and anycast
  - Strategy layer in routers
- Receiver-driven transport
  - Error control performed by receiver
  - Congestion control under research



- Request each video packet
  - Similar to HTTP streaming
  - Difference: request network packets, not chunks
- Receiver-driven stream adaptation *looks* straightforward







• Request video /a/b/c.mp4







- Current rationale: adapt stream quality based on end-to-end bandwidth estimation
  - Packets arrive quickly? Increase quality
  - Packets arrive late? Decrease quality

- Current rationale: adapt stream quality based on end-to-end bandwidth estimation
  - Packets arrive quickly? Increase quality
  - Packets arrive late? Decrease quality
- Hard to estimate *end-to-end bandwidth* in CCN
  - Content source is unknown to receiver

1.  $U_1$  starts with high quality



- 1.  $U_1$  starts with high quality
- 2. Congestion in  $R_2 S_1$ 
  - U<sub>1</sub> switches to low quality



• What does R<sub>1</sub> do?



• What does R<sub>1</sub> do? Forward Interest to R<sub>3</sub>



- What does R<sub>1</sub> do? Forward Interest to R<sub>3</sub>
- What if  $R_1$ - $S_2$  even worse than  $R_1$ - $S_1$ ?



- What does R<sub>1</sub> do? Forward Interest to R<sub>3</sub>
- What if  $R_1$ - $S_2$  even worse than  $R_1$ - $S_1$ ?
  - Client switches back to high?



- Real-time delivery
  - Proactively transmit Interests for upcoming packets
- Native multicast support



- Receiver-driven layered multicast
- Case study: H264 Scalable Video Coding
  - Dependency ID (DID)
  - Quality ID (QID)
  - Temporal ID (TID)
  - Interest: /live-stream/DID;/QID;/TID;/[packet]
- Simple network operation
  - No specific Media Aware Network Elements
  - No multicast JOIN-LEAVE messages



- Packet caches in routers
- Assist in error recovery
  - Cache replacement policy according to packet content
  - I frames > P frames > B frames
  - Video packetization
- Complicates end-to-end bandwidth estimation [3]



[3] Grandl, Su and Westphal, "On the Interaction of Adaptive Video Steaming with Content-Centric Networking," in Packet Video Workshop 2013.

- Overhead caused by Interests
  - One Interest per Data
- Asymmetric /congested uplinks?
- Interest Aggregation [4]
  - Single Interest requests multiple Data packets
  - Additional complexity in routers
  - What if lost?
- Persistent Interests [5]
  - One Interest for all streaming Data packets
  - Similar to IP multicast (channel mode)
  - Longer lifetime than plain Interest
  - PIT size?

[4] Byun, Lee and Jang, "Adaptive Flow Control via Interest Aggregation in CCN,", in IEEE ICC 2013.

[5] Tsilopoulos and Xylomenos, "Supporting Diverse Traffic Types in Information Centric networks," in ACM SIGCOMM ICN workshop 2011.



### **CCN** Summary

	Improved	Unclear	Problematic
Video on Demand	Native anycast support. Enhanced retransmission-based error control with in- network packet-level caching.	End-to-end throughput estimation for stream adaptation.	Network overhead for explicitly requesting individual Data packets
Live Streaming	Enhanced retransmission-based error control with in- network packet-level caching. Packet distinction in caching policies.	Service degradation in asymmetric links. Lost Interests upstream result in missing Data on the downstream.	

#### Publish-Subscribe Internetworking (PSI)

- 3 distinct network functions
  - Rendezvous
  - Topology Management & Path Formation
  - Forwarding
- Decouple routing from forwarding
  - Centralized route selection
  - Explicit-routing, Bloom filter-based
- Pub/sub API
- Abstract notion of content item
  - Not strictly a network packet
  - Could be a larger data unit: chunk or entire file, media stream
- Push-based



#### Publish-Subscribe Internetworking (PSI)

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Operation

- 1. Producer publishes item (announcement)
- 2. Consumer subscribes to item
- 3. Network locates item
- 4. Computes publisher  $\rightarrow$  subscriber path
  - Source route
  - Hands it to publisher
- 5. Publisher transmits data over specified path
  - Sender-driven or receiver-driver transport

### 

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(4)

- 1. Subscribe to video
  - Obtain metadata



- 1. Subscribe to video
  - Obtain metadata
- 2. Subscribe to each piece



- Granularity of video pieces
- Small pieces
  - + Receiver-driven stream adaptation
  - Scalability: number of announcements to Rendezvous
  - Amount of subscriptions: delay for resolution-path formation
- Large pieces
  - Coarse-grained stream adaptation
  - + Less announcements
  - + Fewer subscriptions

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- What we have not seen *yet*:
- Utilize centralized control plane
- Network selects video source and quality on behalf of users
  - QoS parameters
- Need to enrich pub/sub primitives
  - Network must understand data
- Tradeoff general purpose with app specific semantics

video.mp4, low quality

# n yet: trol plane

### Live Streaming in PSI

- Name the stream, not each packet
  - Channel mode, similar to IP multicast
  - + One subscription only
  - No packet caches in routers
- Centralized multicast tree computation
  - Optimization benefits, e.g.
    Steiner trees [6, 7]
  - Increased delays



[6] Li et al., "ESM: Efficient and scalable data center multicast routing," IEEE/ACM Transactions on Networking 2012.[7] Tsilopoulos et al., "Efficient real-time information delivery in future internet publish-subscribe networks," ICNC 2013.

#### **PSI Summary**

	Improved	Unclear	Problematic
Video on Demand	Native anycast support. Optimal path selection through centralized route control.	End-to-end throughput estimation. Optimal path selection requires extensions to pub/sub primitives.	Delays for resolution of subscriptions and unsubscriptions.
Live Streaming	Optimal multicast delivery through centralized route control.	Scalability of centralized multicast tree construction (with dynamic user behavior).	

#### Thank you

**Questions?**